IC637 Program Analysis

Lecture 4: Octagon Domain

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2025 Fall

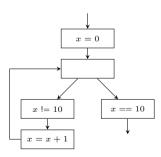
Review: Interval Domain

- Interval Domain: Abstract domain $\hat{\mathbb{Z}} = \{\bot\} \cup \{[l,u] \mid l,u \in \mathbb{Z} \cup \{-\infty,\infty\}, l \leq u\}$
 - Partial order: $\bot \sqsubseteq x$, $[l_1,u_1] \sqsubseteq [l_2,u_2]$ iff $l_2 \le l_1$ and $u_1 \le u_2$
 - Join: $[l_1, u_1] \sqcup [l_2, u_2] = [\min(l_1, l_2), \max(u_1, u_2)]$
 - Meet: $[l_1, u_1] \sqcap [l_2, u_2] = [\max(l_1, l_2), \min(u_1, u_2)]$ if overlap, \perp otherwise
- Worklist Algorithm: Iterative fixed-point computation
 - Widening phase: Apply widening at loop headers until convergence
 - Narrowing phase: Apply narrowing to refine results
- Widening & Narrowing: Essential for termination in infinite domains
 - Widening (∇) : Extrapolates to infinity to ensure convergence
 - Narrowing (\triangle): Refines over-approximations from widening
 - Widening with thresholds: Uses predefined values to improve precision
- **Key Insight**: Balance between precision and scalability (e.g., termination)

Example

- Describe the result of the interval analysis:
 - 1. without widening
 - 2. with widening/narrowing

```
void main() {
   int x = 0;
   while (x != 10) {
        x = x + 1;
   }
}
```

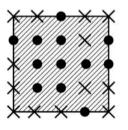


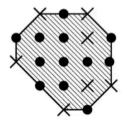
Discussion

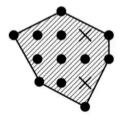
• Give an example program that cannot be precisely analyzed by interval domain.

Relational Abstract Domains

• Intervals vs Octacons vs Polyhedra







The Octagon Domain

• Focus: Core idea of the Octagon domain

```
void main(){
                                           Octagon analysis :
   int a[10];
   x = 0; y = 0;
   while (x < 9) {
x++; y++;
                                            → Interval analysis :
```

• $(N+1) \times (N+1)$ matrix (N: number of variables) e.g.,

$$\begin{array}{c|ccccc} & 0 & x & y \\ \hline 0 & 0 - 0 & x - 0 & y - 0 \\ x & 0 - x & x - x & y - x \\ y & 0 - y & x - y & y - y \end{array}$$

Example

A DBM represents a set of program states (N-dim points)

$$\gamma \left(\begin{bmatrix} 0 & 10 & \infty \\ -1 & 0 & -1 \\ 0 & 1 & 0 \end{bmatrix} \right) = \{ (x,y) \mid 1 \le x \le 10, \ 0 \le y, \ y - x \le -1, \ x - y \le 1 \}$$

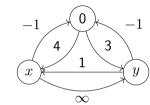
$$\gamma \left(\begin{bmatrix} 0 & 5 & \infty \\ \infty & 0 & -1 \\ 0 & 2 & 0 \end{bmatrix} \right) =$$

- Quiz: find a matrix M such that $\gamma(M) = \emptyset$
- Question: can two different DBMs represent the same set of points?

• A DBM can also be represented by a directed graph

	0	x	y
0	0	4	3
x	-1	0	$+\infty$
y	-1	1	0

$$\iff$$



$$\begin{array}{c|ccccc} & 0 & x & y \\ \hline 0 & 0 & 10 & \infty \\ x & -1 & 0 & -1 \\ y & 0 & 1 & 0 \\ \end{array}$$



Two different DBMs can represent the same set of points

$$\gamma \begin{pmatrix} +\infty & 4 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{pmatrix} = \gamma \begin{pmatrix} 0 & 5 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{pmatrix}$$

• How can we check if two DBMs represent the same set of points?

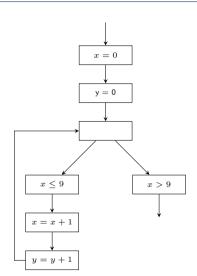
Closure (normalization) via the Floyd-Warshall algorithm

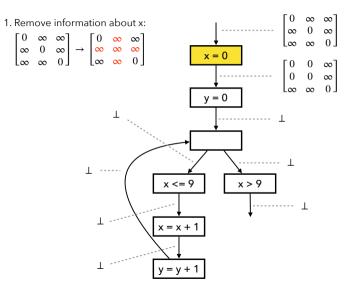
$$\begin{bmatrix} +\infty & 4 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix}^* = \begin{bmatrix} 0 & 4 & 3 \\ -1 & 0 & 2 \\ -1 & 1 & 0 \end{bmatrix}$$
$$\begin{bmatrix} 0 & 5 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix}^* = \begin{bmatrix} 0 & 4 & 3 \\ -1 & 0 & 2 \\ -1 & 1 & 0 \end{bmatrix}$$

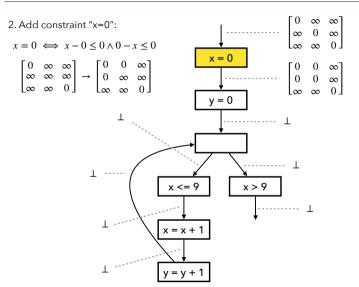
$$\begin{aligned} & \textbf{for } n = 0 \textbf{ to } n - 1 \textbf{ do} \\ & dist[n][n] = 0 \\ & \textbf{for } k = 0 \textbf{ to } n - 1 \textbf{ do} \\ & \textbf{ for } i = 0 \textbf{ to } n - 1 \textbf{ do} \\ & \textbf{ for } j = 0 \textbf{ to } n - 1 \textbf{ do} \\ & dist[i][j] \leftarrow \min(dist[i][j], dist[i][k] + dist[k][j]) \end{aligned}$$

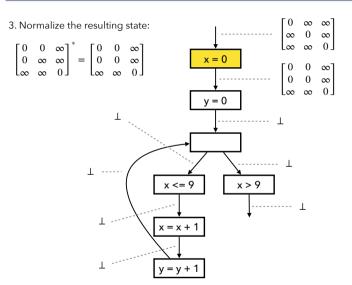
Fixed Point Computation with Widening

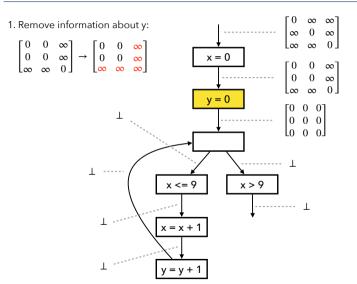
```
int x = 0;
int y = 0;
while (x <= 9) {
    x = x + 1;
    y = y + 1;
}</pre>
```

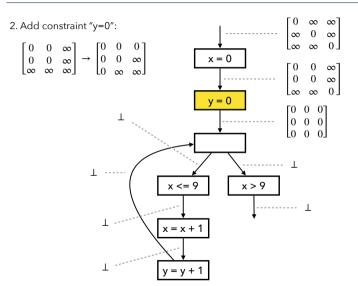


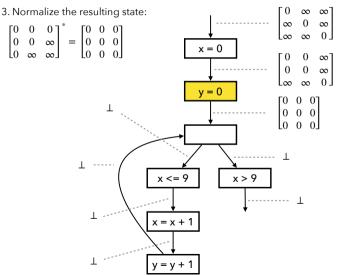


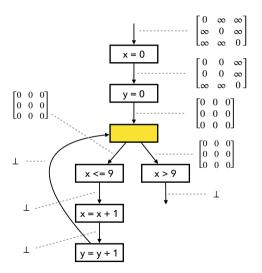


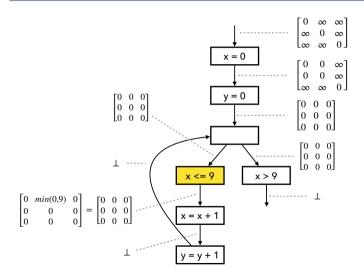


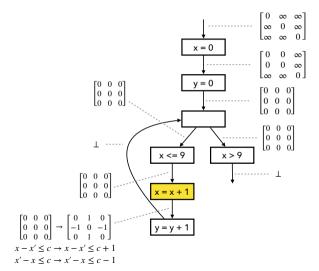


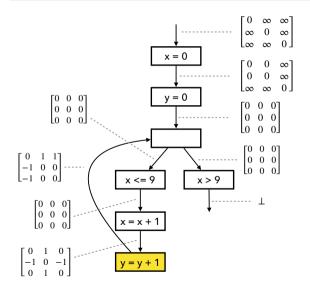


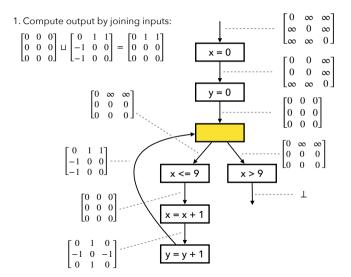


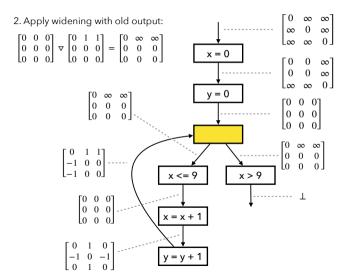


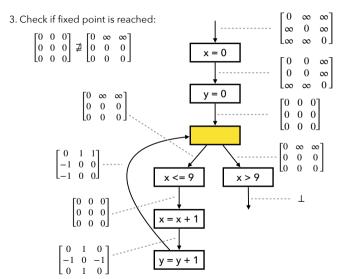


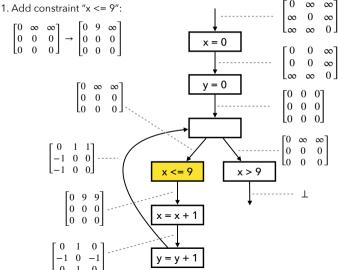


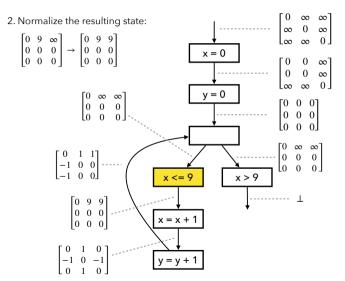


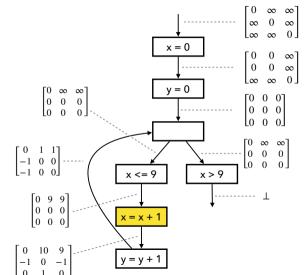




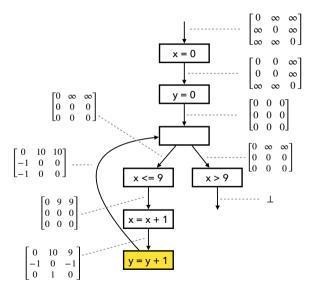


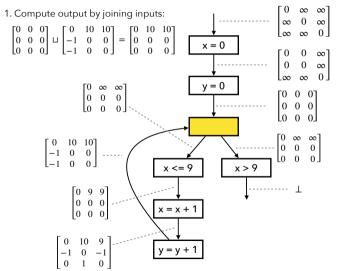


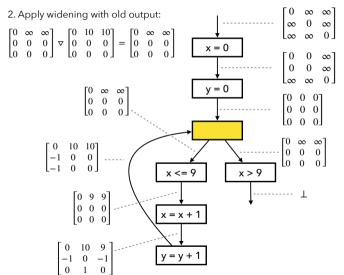


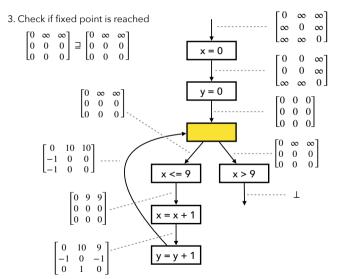


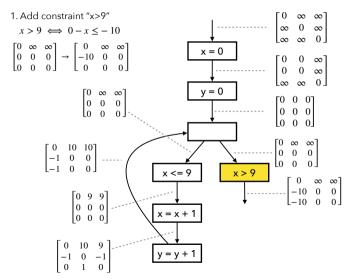
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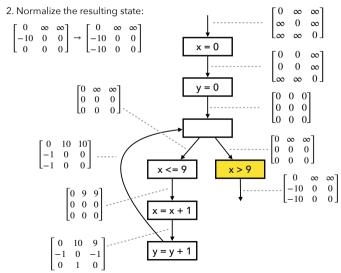


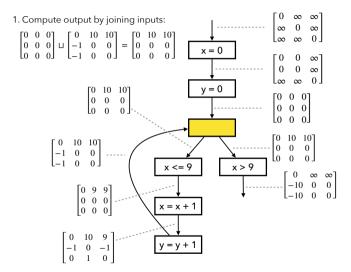


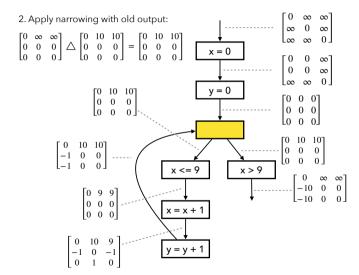


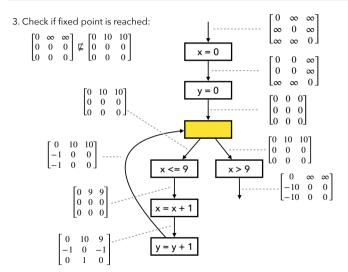


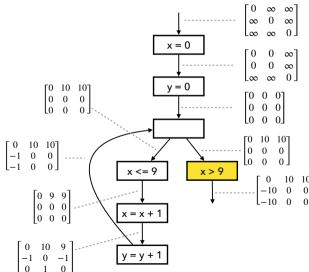






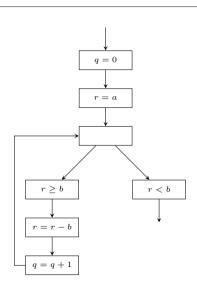






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```
//a >= 0; b >= 0;
int q = 0;
int r = a;
while (r >= b) {
    r = r - b;
    q = q + 1;
}
assert(q >= 0);
assert(r >= 0);
```



Static Analysis Use Cases: Infer

• https://github.com/facebook/infer/





Infer is a static analysis tool for Java, C++, Objective-C, and C. Infer is written in OCaml.

- Running Infer: e.g.,
 - infer capture make
 - infer analyze

Infer's Intermediate Language

https://github.com/facebook/infer/blob/main/infer/src/IR/Sil.mli

```
type instr =
  | Load of {id: Ident.t; e: Exp.t; typ: Typ.t; loc: Location.t}
      (** Load a value from the heap into an identifier.
          [id = *e:tvp] where
          - [e] is an expression denoting a heap address
          - [typ] is the type of [*e] and [id]. *)
  | Store of {e1: Exp.t; typ: Typ.t; e2: Exp.t; loc: Location.t}
      (** Store the value of an expression into the heap.
          [*e1:tvp = e2] where

    [e1] is an expression denoting a heap address

          - [tvp] is the tvpe of [*e1] and [e2]. *)
```

Summary

Static Analysis Principles

- Choose appropriate abstract domains (intervals, octagons, polyhedra)
- Balance precision vs. scalability

Interval Domain

- Simple non-relational domain: [l, u] bounds
- Widening/narrowing for fixed-point computation
- Limited expressiveness for relational properties

Octagon Domain

- Relational domain using Difference Bound Matrices (DBM)
- Constraints: $\pm x_i \pm x_j \le c$
- Floyd-Warshall algorithm for closure/normalization

Practical Tools

Facebook Infer: Industrial static analyzer